

Flexible list colorings: Maximizing the number of requests satisfied

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Precoloring Extension Problem

- Recall proper k -coloring, proper k -list coloring, $\chi(G)$, $\chi_\ell(G)$.

Precoloring Extension Problem

- Is there a proper coloring of a graph G subject to some of the vertices having **prescribed colors**? A function r with non-empty domain $D \subseteq V(G)$ and co-domain of a palette of colors.

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- Is there a proper coloring of a graph G subject to some of the vertices having **prescribed colors**? A function r with non-empty domain $D \subseteq V(G)$ and co-domain of a palette of colors.
- This is a classical problem that has been studied under many contexts. Its not always possible to have such a proper coloring, so typically we look for restrictions on the structure induced by the precolored vertices, etc.

Precoloring Extension Problem

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- What if such a precoloring does not extend, can we instead ask for a coloring which matches the precoloring on many vertices (say, on a constant fraction of the precolored vertices)?

Precoloring Extension Problem

- Is there a proper coloring of a graph G subject to some of the vertices having **prescribed colors**? A function r with non-empty domain $D \subseteq V(G)$ and co-domain of a palette of colors.
- What if such a precoloring does not extend, can we instead ask for a coloring which matches the precoloring on many vertices (say, on a constant fraction of the precolored vertices)?
- **Always possible**, by permuting the colors in a k -coloring of G , we can easily obtain a k -coloring of G that matches r on at least $|dom(r)|/k$ vertices.

List Coloring with Requests

- [Dvořák, Norin, and Postle \(2019\)](#): A proper list coloring, but a preferred color is given for some subset of vertices and we wish to color as many vertices in this subset with its preferred color as possible; a flexible version of the classical precoloring extension problem.

List Coloring with Requests

- Given a graph G and a list assignment L of G .
A **request** of L is a function r with non-empty domain $D \subseteq V(G)$ such that $r(v) \in L(v)$ for each $v \in D$.
For any $\epsilon \in (0, 1]$, (G, L, r) is **ϵ -satisfiable** if there exists a proper L -coloring f of G such that $f(v) = r(v)$ for at least $\epsilon|D|$ vertices in D .

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- (G, L) is **ϵ -flexible** if (G, L, r) is ϵ -satisfiable whenever r is a request of L .
 G is **(k, ϵ) -flexible** if (G, L) is ϵ -flexible whenever L is a k -assignment for G .

List Coloring with Requests

- (G, L) is ϵ -flexible if (G, L, r) is ϵ -satisfiable whenever r is a request of L .
 G is (k, ϵ) -flexible if (G, L) is ϵ -flexible whenever L is a k -assignment for G .
- If G is (k, ϵ) -flexible, then it immediately follows:
 - (i) G is (k', ϵ') -flexible for any $k' \geq k$ and $\epsilon' \leq \epsilon$;
 - (ii) any spanning subgraph H of G is (k, ϵ) -flexible;
 - (iii) G is k -choosable.

Previous works

- Dvořák, Norin, and Postle mostly focused on d -degenerate graphs, which are known to be $(d + 1)$ -choosable. They showed that d -degenerate graphs are all $(d + 2, \epsilon)$ -flexible for some $\epsilon > 0$.
- The main open problem Dvořák-Norin-Postle asked was:
Are all d -degenerate graphs, $(d + 1, \epsilon(d))$ -flexible for some $\epsilon(d) > 0$?

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Are all d -degenerate graphs, $(d + 1, \epsilon(d))$ -flexible for some $\epsilon(d) > 0$?
- They showed there exists an $\epsilon > 0$ such that every planar graph G is $(6, \epsilon)$ -flexible.
- Flexible list coloring has been studied extensively for planar graphs that are 5-choosable, and for restricted subclasses of planar graphs that are k -choosable with $k < 5$.
Several papers on (k, ϵ) -flexibility, with $k \in \{5, 4, 3\}$, of planar graphs with large enough girth or excluding certain cycles.

Motivation for our work

- Find the largest possible ϵ for which G is (k, ϵ) -flexible; that is, one would prefer to have a larger portion of the requested colors on vertices satisfied.

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- Find the largest possible ϵ for which G is (k, ϵ) -flexible; that is, one would prefer to have a larger portion of the requested colors on vertices satisfied.
- Only previous result of this flavor is the following.

Bradshaw, T. Masařík, L. Stacho (2022): Let G be a connected graph with $\Delta(G) \geq 3$ that is not a copy of $K_{\Delta(G)+1}$. Then, G is $(\Delta(G), 1/(6\Delta(G)))$ -flexible. Moreover, $1/(6\Delta(G))$ is within a constant factor of being best possible.

Improving older results - I

Theorem (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose G is d -degenerate. Then, G is $(d + 2, \frac{1}{2^{d+1}})$ -flexible.

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Suppose G is d -degenerate. Then, G is $(d + 2, \frac{1}{2^{d+1}})$ -flexible.

- Proved using a randomized algorithm:

We can order vertices of G as v_1, \dots, v_n such that for all $i \in [n]$, v_i has at most d neighbors v_j with $j > i$, and order each list of colors, $L(v_i)$, such that the requested color (if it exists) is the first color. Now pick uniformly from the first two available colors as we color the vertices in order.

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Theorem (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose G is d -degenerate. Then, G is $(d + 2, \frac{1}{2^{d+1}})$ -flexible.

- This improves the following:

Theorem (Dvořák, Norin, and Postle (2019))

Suppose G is d -degenerate. Then,

G is $\left(d + 2, \frac{1}{(d + 2)^{(d+1)^2}}\right)$ -flexible.

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Theorem (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose G is d -degenerate. Then, G is $(d + 2, \frac{1}{2^{d+1}})$ -flexible.

Theorem (K., Mathew, Mudrock, Pelsmayer (2022+))

Let G be an s -choosable graph. Then, G is $(s + 1, 1/\chi(G^2))$ -flexible.

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Let G be an s -choosable graph. Then, G is $(s + 1, 1/\chi(G^2))$ -flexible.

This yields an improvement on our first theorem above for d -degenerate graphs with maximum degree $\Delta < 2^d/d$.

Corollary (K., Mathew, Mudrock, Pelsmayer (2022+))

Let G be a d -degenerate graph with maximum degree Δ . If G is s -choosable, then G is $(s + 1, 1/(\Delta(2d - 1) + d - d^2 + 1))$ -flexible.

Improving older results - II

- If we focus solely on k and allow arbitrarily small $\epsilon > 0$, then our colorings need only satisfy the color request at a single vertex. Then, without loss of generality, we need to only study requests with domain of size 1, as those have the most restrictive requirement.

Improving older results - II

- If we focus solely on k and allow arbitrarily small $\epsilon > 0$, then our colorings need only satisfy the color request at a single vertex. Then, without loss of generality, we need to only study requests with domain of size 1, as those have the most restrictive requirement.
- Dvořák, Norin, and Postle say “A *necessary condition for flexibility is that requests with singleton domain can be satisfied. Coming back to the case of d -degenerate graphs with lists of size $d + 1$, even proving this necessary condition is non-trivial and we can only do it in the special case that $d + 1$ is a prime.*”

Theorem (Dvořák, Norin, and Postle (2019))

Let $d \geq 2$ such that $d + 1$ is a prime. If G is a d -degenerate graph, L is a $(d + 1)$ -assignment, and r is a request for G with domain of size 1, then (G, L, r) is 1-satisfiable.

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Theorem (Dvořák, Norin, and Postle (2019))

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Using the Alon-Tarsi Theorem we are able to extend their Theorem to all d for bipartite d -degenerate graphs.

Theorem (K., Mathew, Mudrock, Pelsmayer (2022+))

For any bipartite d -degenerate graph G with a $(d + 1)$ -list assignment L and request r with domain D of size 1, (G, L, r) is 1-satisfiable.

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Theorem (K., Mathew, Mudrock, Pelsmayer (2022+))

For any bipartite d -degenerate graph G with a $(d + 1)$ -list assignment L and request r with domain D of size 1, (G, L, r) is 1-satisfiable.

We can order vertices of G as v_1, \dots, v_n such that for all $i \in [n]$, v_i has at most d neighbors v_j with $j > i$. Suppose that v_k is the vertex in the domain of r .

We show that there is a way, using Menger's Theorem, to orient the edges of G to obtain a digraph in which every vertex has out-degree at most d and v_k has out-degree zero. The result then follows from Alon-Tarsi.

Maximizing the number of vertex requests satisfied

For each graph G , what is the largest ϵ so that G is (k, ϵ) -flexible for some k ?

Maximizing the number of vertex requests satisfied

- It is possible that $r(v)$ is the same color for all $v \in D$; for example, let L be the k -list assignment such that $L(v) = [k]$ for all $v \in V(G)$ and let $r(v) = 1$ for all $v \in D$. Then at most $\alpha(G[D])$ vertices in D will have their request fulfilled. So, $\epsilon \leq \min_{\emptyset \neq D \subseteq V(G)} \alpha(G[D]) / |D|$ for any (k, ϵ) -flexible graph G .

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- The Hall ratio of a graph G is $\rho(G) = \max_{\emptyset \neq H \subseteq G} \frac{|V(H)|}{\alpha(H)}$.

The Hall ratio was first studied in 1990 by Hilton and Johnson Jr. under the name Hall-condition number in the context of list coloring. In the past 30 years, the Hall ratio has received much attention due to its connection with both list and fractional coloring.

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Proposition (K., Mathew, Mudrock, Pelsmayer (2022+))

There exists k such that G is (k, ϵ) -flexible if and only if $\epsilon \leq 1/\rho(G)$.

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- We define the **list flexibility number** of G , denoted $\chi_{flex}(G)$, to be the smallest k such that G is $(k, 1/\rho(G))$ -flexible.

Maximizing the number of vertex requests satisfied

- We define the **list flexibility number** of G , denoted $\chi_{flex}(G)$, to be the smallest k such that G is $(k, 1/\rho(G))$ -flexible.

List flexibility number

- $\chi_{flex}(G)$ is the smallest k such that G is $(k, 1/\rho(G))$ -flexible.
- Is maximizing ϵ meaningfully different from a flexible list coloring with smaller $\epsilon > 0$? In other words, **are there any graphs G and $k \in \mathbb{N}$ such that G is (k, ϵ) -flexible for some $\epsilon > 0$, but G is not $(k, 1/\rho(G))$ -flexible?** The following result shows that the answer is yes.

Proposition (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose $G = K_{3,7}$. Then, G is $(3, 1/10)$ -flexible and $\chi_{flex}(G) > 3$.

List flexibility number

- $\chi_{flex}(G)$ is the smallest k such that G is $(k, 1/\rho(G))$ -flexible.
- $\chi_{\ell}(G) \leq \chi_{flex}(G) \leq \Delta(G) + 1$.
- It follows that $\chi_{flex}(K_n) = n$ and $\chi_{flex}(C_k) = 3$ for odd k . It is natural to ask whether a Brooks-type theorem is true for χ_{flex} as well.
Question: What are all the graphs G such that $\chi_{flex}(G) = \Delta(G) + 1$?

χ_{flex} vs degeneracy

- Can the Dvořák-Norin-Postle-Conjecture that d -degenerate graphs G are $(d + 1, \epsilon(d))$ -flexible for some $\epsilon(d) > 0$ be strengthened to $\chi_{flex}(G) \leq d + 1$?

Question: Does there exist a graph G with degeneracy d satisfying $\chi_{flex}(G) > d + 1$?

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Question: Does there exist a graph G with degeneracy d satisfying $\chi_{flex}(G) > d + 1$?

- Alon (2000) showed for any graph G with degeneracy d , $(1/2 - o(1)) \log_2(d + 1) \leq \chi_{\ell}(G)$ and is sharp up to a factor of 2. How sharp is this lower bound for $\chi_{flex}(G)$?

Question: Suppose

$\mathcal{F}(d) = \min\{\chi_{flex}(G) : \text{the degeneracy of } G \text{ is at least } d\}$.
What is the asymptotic behavior of $\mathcal{F}(d)$ as $d \rightarrow \infty$?

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What is the asymptotic behavior of $\mathcal{F}(d)$ as $d \rightarrow \infty$?

- We are able to show $\mathcal{F}(d) = O(d)$ while Alon's result implies $\mathcal{F}(d) = \Omega(\log_2(d))$, as $d \rightarrow \infty$.

χ_{flex} vs degeneracy

Theorem (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose G is an n -vertex, s -choosable graph with $\chi_{flex}(G) = m$ and $n \geq 2$. Let $l = \lceil n/\rho(G) \rceil$. Let $J = G \vee G$. Then, for any real number $r > 2$,

$$\chi_{flex}(J) \leq \max \left\{ \left\lceil l + \frac{\log_2(2n-l)}{1-H(1/r)} \right\rceil, \lceil r(s-1) + l \rceil, m \right\}.$$

- H is the binary Entropy function.

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Theorem (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose n is a positive even integer, and $G = P_n \vee P_n$. If $p \in (0, 1)$ and $k > \max\{2, n/2\}$ satisfy $(n/2)(1-p)^{k-2} + np^{k-1-n/2}(p + (k - n/2)(1-p)) \leq 1$, then $\chi_{flex}(G) \leq k$.

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Corollary

Suppose n is a positive even integer satisfying $n \geq 50$. Then, $\chi_{flex}(P_n \vee P_n) \leq \lceil n/2 + \ln(n) \rceil$.

Since degeneracy of $P_n \vee P_n$ is $n + 1$, it follows

Corollary

$\mathcal{F}(d)$ grows no faster than $d/2$ as $d \rightarrow \infty$.

$\chi_{flex}(G)$ vs $\chi_{\ell}(G)$

- Recall $\chi_{\ell}(G) \leq \chi_{flex}(G) \leq \Delta(G) + 1$.
- It is natural to ask whether $\chi_{flex}(G)$ can be bounded above by a function of $\chi_{\ell}(G)$.
Question: Does there exist a function f such that for every graph G , $\chi_{flex}(G) \leq f(\chi_{\ell}(G))$?

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- It is natural to ask whether $\chi_{flex}(G)$ can be bounded above by a function of $\chi_{\ell}(G)$.
Question: Does there exist a function f such that for every graph G , $\chi_{flex}(G) \leq f(\chi_{\ell}(G))$?
- We show that there is no universal constant C such that $\chi_{flex}(G) \leq \chi_{\ell}(G) + C$.

$\chi_{flex}(G)$ vs $\chi_{\ell}(G)$

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Proposition (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose that $l \in \mathbb{N}$ and $t_0 = \sum_{i=0}^l \binom{2l+1}{2l+1-i} (2l)^i$.

(i) For $t \geq t_0$, $\chi_{flex}(K_{2l+1,t}) \geq 2l + 2$.

(ii) For $s \leq t_0$, $\chi_{\ell}(K_{2l+1,s}) \leq \lceil 3l/2 \rceil$ whenever $l \geq 100000$.

Corollary

Let $t_0 = \sum_{i=0}^l \binom{2l+1}{2l+1-i} (2l)^i$.

For each $l \geq 100000$, $\chi_{flex}(K_{2l+1,t}) \geq \frac{4}{3}\chi_{\ell}(K_{2l+1,t}) + \frac{4}{3}$.

χ_{flex} vs List packing

- An idea implicit in an earlier work (for k -trees):

Proposition (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose that G is a graph, L is a k -assignment for G , and there is a set S of mk proper L -colorings such that for each vertex $v \in V(G)$ and each color $c \in L(v)$, vertex v is colored by c in exactly m of the L -colorings of S . Then, (G, L, r) is $1/k$ -satisfiable for any request r of L .

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- For a non-trivial tree T , it is easy to see that for any 2-assignment L , there exist 2 proper L -colorings that are distinct on each vertex. Here $m = 1$ and $k = 2$. Since $\rho(T) = 2$ and $\chi_{flex}(T) \geq \chi_\ell(T) = 2$, above Proposition implies $\chi_{flex}(T) = 2$.

χ_{flex} vs List packing

- **List packing** is a relatively new notion that was first suggested by [Alon, Fellows, and Hare \(1996\)](#), and formally defined in a recent paper of [Cambie, Batenburg, Davies, Kang \(2021+\)](#) to appear in [RSA 2023](#).

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- Let L be a list assignment for a graph G . An **L -packing of G of size k** is a set of proper L -colorings $\{f_1, \dots, f_k\}$ of G such that $f_i(v) \neq f_j(v)$ whenever $i, j \in [k]$, $i \neq j$, and $v \in V(G)$. The **list packing number** of G , denoted $\chi_\ell^*(G)$, is the least k such that G has a proper L -packing of size k whenever L is a k -assignment for G .
- $\chi(G) \leq \chi_\ell(G) \leq \chi_\ell^*(G)$.

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Proposition (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose that G is a graph, L is a k -assignment for G , and there is a set \mathcal{S} of mk proper L -colorings such that for each vertex $v \in V(G)$ and each color $c \in L(v)$, vertex v is colored by c in exactly m of the L -colorings of \mathcal{S} . Then, (G, L, r) is $1/k$ -satisfiable for any request r of L .

Corollary (K., Mathew, Mudrock, Pelsmayer (2022+))

For any graph G , G is $(\chi_\ell^(G), 1/\chi_\ell^*(G))$ -flexible.*

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- The **list packing number** of G , denoted $\chi_\ell^*(G)$, is the least k such that G has a proper L -packing of size k whenever L is a k -assignment for G .

$$\chi(G) \leq \chi_\ell(G) \leq \chi_\ell^*(G).$$

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- The **list packing number** of G , denoted $\chi_{\ell}^*(G)$, is the least k such that G has a proper L -packing of size k whenever L is a k -assignment for G .

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Corollary (K., Mathew, Mudrock, Pelsmayer (2022+))

For any graph G , G is $(\chi_{\ell}^(G), 1/\chi_{\ell}^*(G))$ -flexible.*

- In view of the Corollary above, and since $\chi_{\ell}(G) \leq \chi_{flex}(G)$, it is natural to ask whether $\chi_{flex}(G)$ can be bounded above by a function of $\chi_{\ell}^*(G)$. More ambitiously,

Conjecture: For any graph G , $\chi_{flex}(G) \leq \chi_{\ell}^*(G)$.

χ_{flex} vs List packing

- Conjecture: For any graph G , $\chi_{flex}(G) \leq \chi_{\ell}^*(G)$.
- Cambie and Härmäläinen (2023+): $\chi_{\ell}^*(K_{3,7}) = 3$,
disprove this conjecture with help of
K., Mathew, Mudrock, Pelsmayer (2022+): $\chi_{flex}(K_{3,7}) > 3$,
and conjecture there are infinitely many counterexamples.

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disprove this conjecture with help of
K., Mathew, Mudrock, Pelsmayer (2022+): $\chi_{flex}(K_{3,7}) > 3$,
and conjecture there are infinitely many counterexamples.
- However, it follows from our Proposition and a result of
Cambie, Batenburg, Davies, Kang (2021+) that:
Every graph G on n vertices is
 $(\chi_{\ell}^*(G), 1/((5 + o(1))\rho(G)(\log n)^2))$ -flexible where the $o(1)$
term tends to 0 as n tends to infinity.

χ_{flex} using List packing

Proposition (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose that G is a graph, L is a k -assignment for G , and there is a set S of mk proper L -colorings such that for each vertex $v \in V(G)$ and each color $c \in L(v)$, vertex v is colored by c in exactly m of the L -colorings of S . Then, (G, L, r) is $1/k$ -satisfiable for any request r of L .

Corollary (K., Mathew, Mudrock, Pelsmayer (2022+))

For any graph G , G is $(\chi_{\ell}^(G), 1/\chi_{\ell}^*(G))$ -flexible.*

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Corollary (K., Mathew, Mudrock, Pelsmayer (2022+))

For any graph G , G is $(\chi_{\ell}^(G), 1/\chi_{\ell}^*(G))$ -flexible.*

- For a non-trivial tree T , it is easy to see that $(\chi_{\ell}^*(T) = 2$. Since $\rho(T) = 2$ and $\chi_{flex}(T) \geq \chi_{\ell}(T) = 2$, above Proposition implies $\chi_{flex}(T) = 2$.
- It follows from above Proposition and a result of [Cambie, Batenburg, Davies, Kang \(2021+\)](#) that:
Every graph G on n vertices is $(\chi_{\ell}^*(G), 1/((5 + o(1))\rho(G)(\log n)^2))$ -flexible where the $o(1)$ term tends to 0 as n tends to infinity.

χ_{flex} using List packing

- Using $\chi_{\ell}^*(P_n) = 2$ and our Proposition, we get

Proposition (K., Mathew, Mudrock, Pelsmayer (2022+))

The grid $P_n \square P_m$ is $(3, 1/3)$ -flexible.

And, we are able to obtain the best possible result for the n -ladder $P_2 \square P_n$.

Proposition (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose $G = P_2 \square P_n$ with $n \geq 2$. Then, G is $(3, 1/2)$ -flexible. Consequently, $\chi_{flex}(G) = 3$.

and, more generally

Proposition (K., Mathew, Mudrock, Pelsmayer (2022+))

Suppose G is (k, ϵ) -flexible. Then $G \square H$ is $(\max\{k, \Delta(H) + \chi_{\ell}(G)\}, \epsilon/\chi(H))$ -flexible.

Thank You! Questions?

- Dvořák, Norin, and Postle (2019): Are all d -degenerate graphs, $(d + 1, \epsilon(d))$ -flexible for some $\epsilon(d) > 0$?
- For any d -degenerate graph G with a $(d + 1)$ -list assignment L and request r with domain D of size 1, show that (G, L, r) is 1-satisfiable.
- What are all the graphs G such that $\chi_{flex}(G) = \Delta(G) + 1$?
- Does there exist a graph G with degeneracy d satisfying $\chi_{flex}(G) > d + 1$?
- What is the asymptotic behavior of $\mathcal{F}(d) = \min\{\chi_{flex}(G) : \text{the degeneracy of } G \text{ is at least } d\}$ as $d \rightarrow \infty$? Linear or Logarithmic in d ?
- Can $\chi_{flex}(G)$ be bounded above by a function of $\chi_{\ell}(G)$?
- Can $\chi_{flex}(G)$ be bounded above by a function of $\chi_{\ell}^*(G)$?
Are there infinitely many counterexamples to $\chi_{flex}(G) \leq \chi_{\ell}^*(G)$?
- Let ϵ_G be the function that maps each $k \in \mathbb{N}$ to the largest ϵ such that G is (k, ϵ) -flexible. Clearly, $\epsilon_G(k) = a/b$ for some integers $0 \leq a \leq b \leq |V(G)|$ and $\epsilon_G(k) \leq 1/\rho(G)$. Study ϵ_G for various G .

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