



New Results on Graph Packing

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Introduction

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We may assume $|V_1| = |V_2| = n$ by adding isolated vertices.

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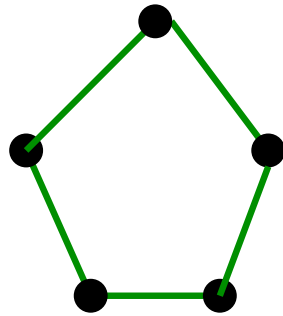
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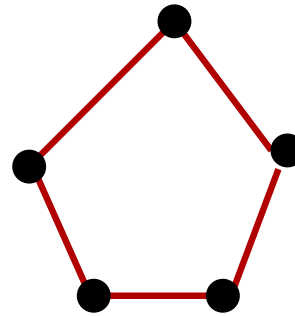
such that the images of the edge sets do not intersect.

- there exists a bijection $V_1 \leftrightarrow V_2$ such that $e \in E_1 \Rightarrow e \notin E_2$.
- G_1 is a subgraph of $\overline{G_2}$.

Examples and Non-Examples

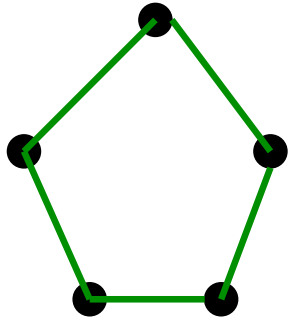


C_5

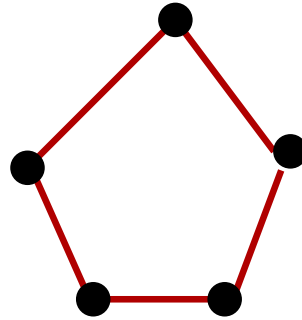


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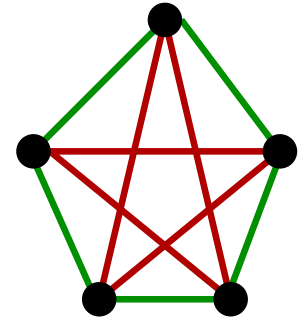
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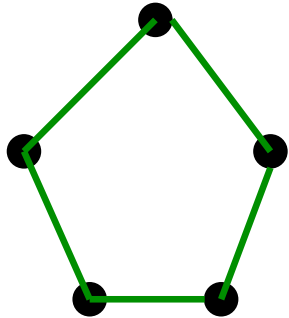


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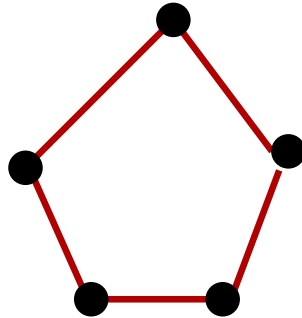


Packing

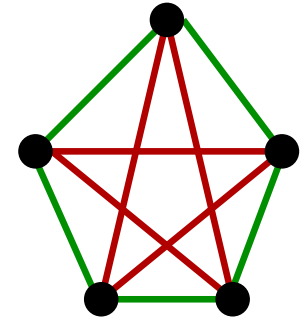
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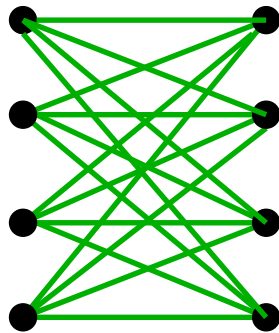
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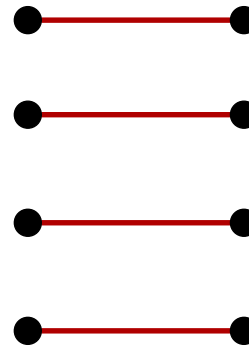
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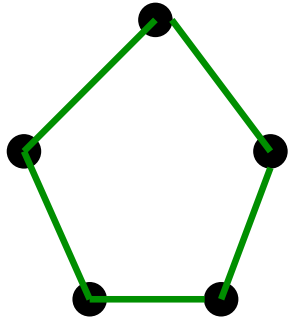


$K_{4,4}$

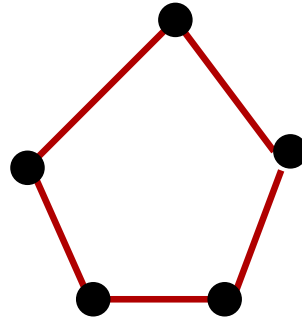


$4 K_2$

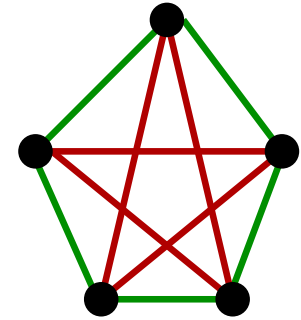
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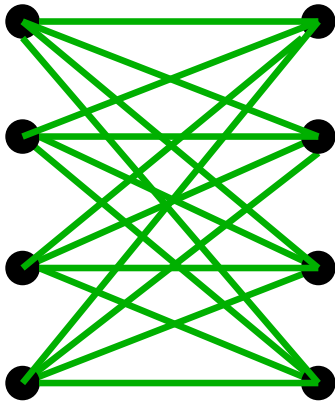
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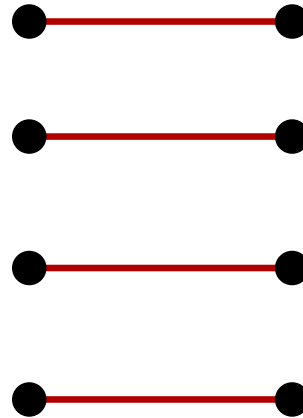
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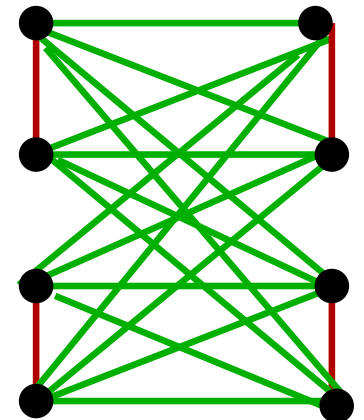
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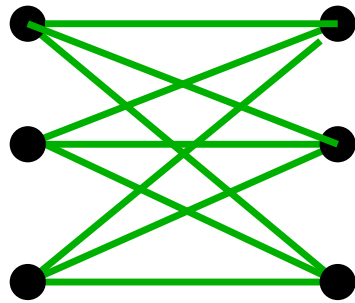


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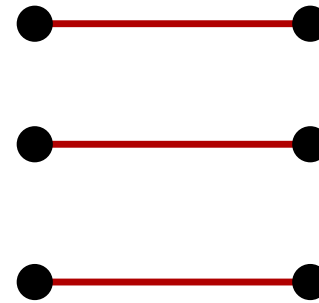


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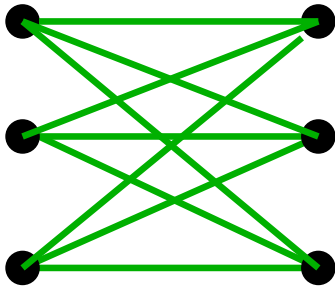


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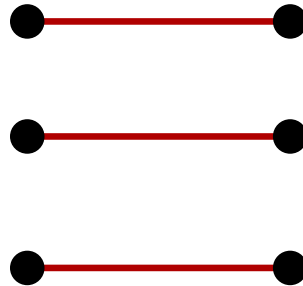


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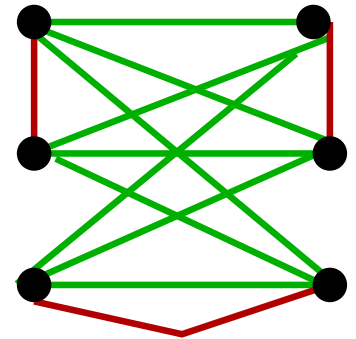
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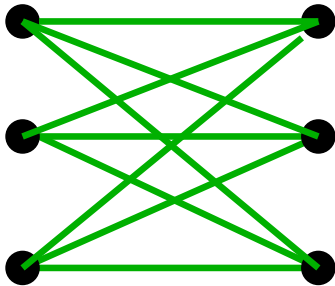


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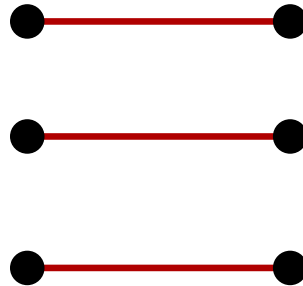


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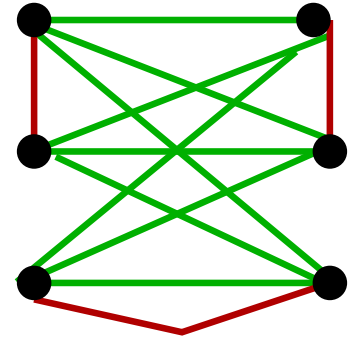
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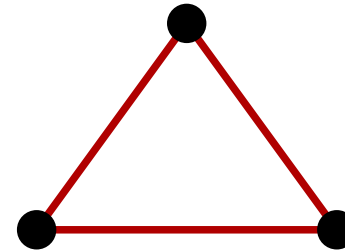
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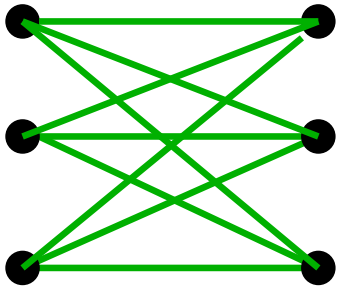


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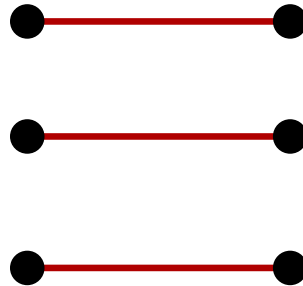


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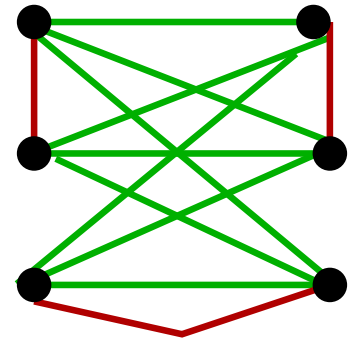
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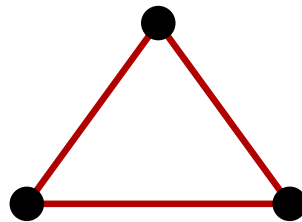
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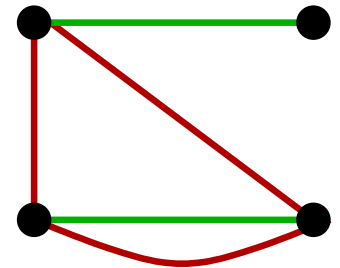
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A Common Generalization

- Hamiltonian Cycle in graph G : Whether the n -cycle C_n packs with \overline{G} .

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- Turan-type problems (forbidden subgraphs).
- Ramsey-type problems.
- “most” problems in Extremal Graph Theory.

A Distinction

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Erdős-Sos Conjecture : Let G be a graph of order n and T be a tree of size k .

If $e(G) < \frac{1}{2}n(n - k)$ then T and G pack.

Sauer and Spencer's Packing Theorem

Theorem [Sauer + Spencer, 1978] :
If $2\Delta_1\Delta_2 < n$, then G_1 and G_2 pack.

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This is sharp.

For n even.

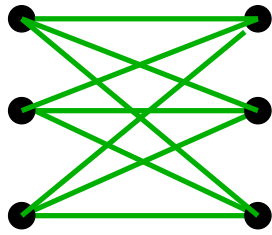
$G_1 = \frac{n}{2}K_2$, a perfect matching on n vertices.

$G_2 \supseteq K_{\frac{n}{2}+1}$, or

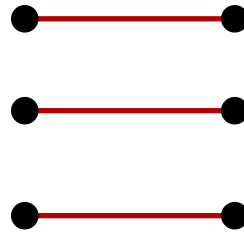
$G_2 = K_{\frac{n}{2}, \frac{n}{2}}$ with $\frac{n}{2}$ odd.

Then, $2\Delta_1\Delta_2 = n$, and G_1 and G_2 do not pack.

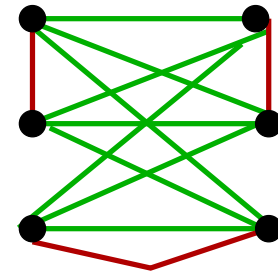
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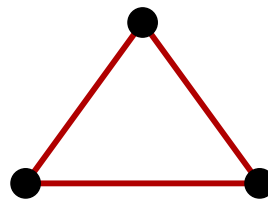
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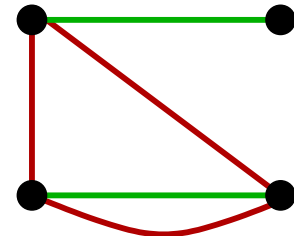
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K_3



No Packing

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Extending the Sauer-Spencer Theorem

Theorem 1 [Kaul + Kostochka, 2005]:

If $2\Delta_1\Delta_2 \leq n$, then

G_1 and G_2 do not pack if and only if one of G_1 and G_2 is a perfect matching and the other either is $K_{\frac{n}{2}, \frac{n}{2}}$ with $\frac{n}{2}$ odd or contains $K_{\frac{n}{2}+1}$.

This result characterizes the extremal graphs for the Sauer-Spencer Theorem.

To appear in *Combinatorics, Probability and Computing*.

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This result can also be thought of as a small step towards the well-known Bollobás-Eldridge conjecture.

Bollobás-Eldridge Graph Packing Conjecture :

If $(\Delta_1 + 1)(\Delta_2 + 1) \leq n + 1$ then G_1 and G_2 pack.

Bollobás-Eldridge Conjecture

Bollobás-Eldridge Graph Packing Conjecture [1978] :
If $(\Delta_1 + 1)(\Delta_2 + 1) \leq n + 1$ then G_1 and G_2 pack.

If true, this conjecture would be sharp, and would be a considerable extension of the Hajnal-Szemerédi theorem on equitable colorings.

The conjecture has only been proved when

$\Delta_1 \leq 2$ [Aigner + Brandt (1993), and Alon + Fischer (1996)], or

$\Delta_1 = 3$ and n is huge [Csaba + Shokoufandeh + Szemerédi (2003)].

Reformulating the Conjecture

Let us consider a refinement of the Bollobás-Eldridge Conjecture.

Conjecture : For a fixed $0 \leq \epsilon \leq 1$.

If $(\Delta_1 + 1)(\Delta_2 + 1) \leq \frac{n}{2}(1 + \epsilon) + 1$, then G_1 and G_2 pack.

For $\epsilon = 0$, this is essentially the Sauer-Spencer Theorem, while $\epsilon = 1$ gives the Bollobás-Eldridge conjecture.

For any $\epsilon > 0$ this would improve the Sauer-Spencer result (in a different way than **Theorem 1**).

Towards the Bollobás-Eldridge Conjecture

Theorem 2 [Kaul + Kostochka + Yu, 2005+]:

For $\epsilon = 0.2$, and $\Delta_1, \Delta_2 \geq 400$,

If $(\Delta_1 + 1)(\Delta_2 + 1) \leq \frac{n}{2}(1 + \epsilon) + 1$, then G_1 and G_2 pack.

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In other words,

Theorem 2 [Kaul + Kostochka + Yu, 2005+]:

For $\Delta_1, \Delta_2 \geq 400$,

If $(\Delta_1 + 1)(\Delta_2 + 1) \leq (0.6)n + 1$, then G_1 and G_2 pack.

This is work in progress.

Some Proof Ideas for Theorem 1

Theorem 1 [Kaul + Kostochka, 2005]:

If $2\Delta_1\Delta_2 \leq n$, then

G_1 and G_2 do not pack if and only if one of G_1 and G_2 is a perfect matching and the other either is $K_{\frac{n}{2}, \frac{n}{2}}$ with $\frac{n}{2}$ odd or contains $K_{\frac{n}{2}+1}$.

We have to analyze the ‘minimal’ graphs that do not pack (under the condition $2\Delta_1\Delta_2 \leq n$).

(G_1, G_2) is a *critical pair* if G_1 and G_2 do not pack, but for each $e_1 \in E(G_1)$, $G_1 - e_1$ and G_2 pack, and for each $e_2 \in E(G_2)$, G_1 and $G_2 - e_2$ pack.

Some Proof Ideas for Theorem 1

Each bijection $f : V_1 \rightarrow V_2$ generates a (multi)graph G_f , with

$$V(G_f) = \{(u, f(u)) : u \in V_1\}$$

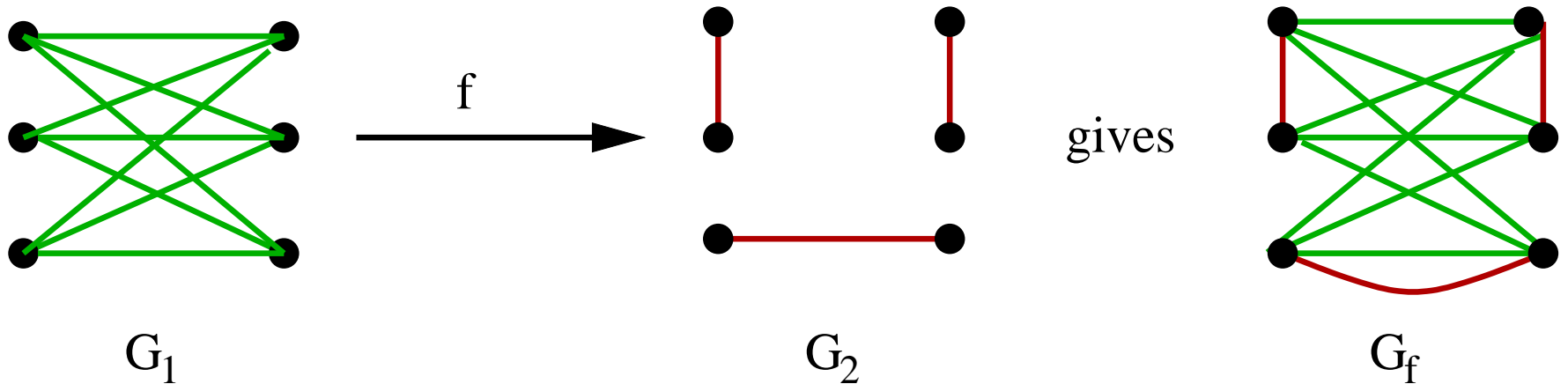
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$$f'(u) = \begin{cases} f(u) & , \quad u \neq u_1, u_2 \\ f(u_2) & , \quad u = u_1 \\ f(u_1) & , \quad u = u_2 \end{cases}$$

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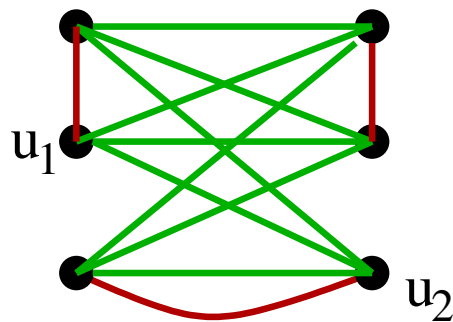
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2-neighbors of u_1 \longleftrightarrow 2-neighbors of u_2

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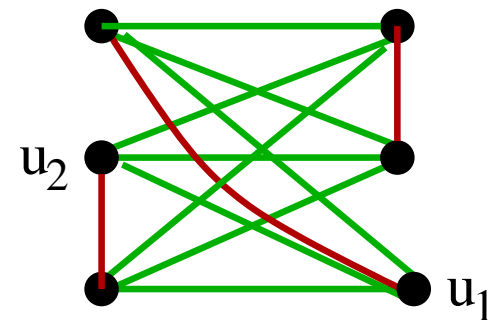
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G_f



(u_1, u_2) -switch



G_f'

Some Proof Ideas for Theorem 1

An important structure that we utilize in our proof is -

$(u_1, u_2; 1, 2)$ -*link* is a path of length two (in G_f) from u_1 to u_2 whose first edge is in E_1 and the second edge is in E_2 .

For $e \in E_1$, an e -*packing* (*quasi-packing*) of (G_1, G_2) is a bijection f between V_1 and V_2 such that e is the only edge in E_1 that shares its incident vertices with an edge from E_2 .

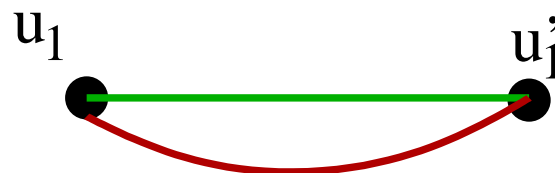
Such a packing exists for every edge e in a critical pair.

Outline of the Proof of Theorem 1

The main tool –

Lemma 1 : Let (G_1, G_2) be a critical pair and $2\Delta_1\Delta_2 \leq n$. Given any $e \in E_1$, in a e -packing of (G_1, G_2) with $e = u_1u'_1$, the following statements are true.

- (i) For every $u \neq u'_1$, there exists either a unique $(u_1, u; 1, 2)$ -link or a unique $(u_1, u; 2, 1)$ -link,
- (ii) there is no $(u_1, u'_1; 1, 2)$ -link or $(u_1, u'_1; 2, 1)$ -link,
- (iii) $2\deg_{G_1}(u_1)\deg_{G_2}(u_1) = n$.

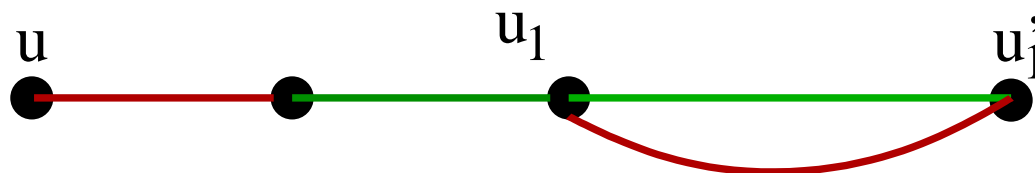


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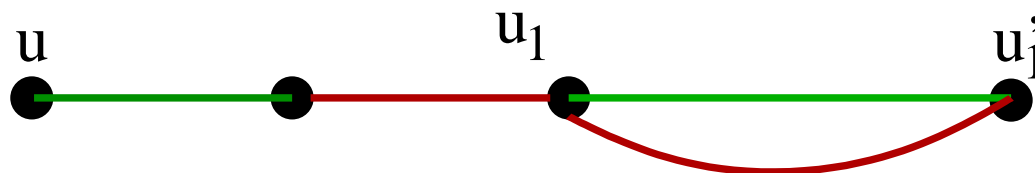


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Outline of the Proof of Theorem 1

Lemma 2 : If $2\Delta_1\Delta_2 = n$ and (G_1, G_2) is a critical pair, then every component of G_i is either K_{Δ_i, Δ_i} with Δ_i odd, or an isolated vertex, or K_{Δ_i+1} , $i = 1, 2$.

Lemma 2 allows us to settle the case of : Δ_1 or $\Delta_2 = 1$.

Then, we have to give a packing for all remaining pairs of graphs, to eliminate their possibility.

Outline of the Proof of Theorem 1

The following Lemma limits the possible remaining pairs of graphs.

Lemma 3 : Let $\Delta_1, \Delta_2 > 1$ and $2\Delta_1\Delta_2 = n$. If (G_1, G_2) is a critical pair and the conflicted edge in a quasi-packing belongs to a component H of G_2 isomorphic to K_{Δ_2, Δ_2} , then every component of G_1 sharing vertices with H is K_{Δ_1, Δ_1} .

Now, we completely eliminate such graphs.

Lemma 4 : Suppose that $\Delta_1, \Delta_2 \geq 3$ and odd, and $2\Delta_1\Delta_2 = n$. If G_1 consists of Δ_2 copies of K_{Δ_1, Δ_1} and G_2 consists of Δ_1 copies of K_{Δ_2, Δ_2} , then G_1 and G_2 pack.

Outline of the Proof of Theorem 1

Now, let's eliminate the only remaining possibility.

Lemma 5 : Let $\Delta_1, \Delta_2 > 1$ and $2\Delta_1\Delta_2 = n$. If every non-trivial component of G_i is K_{Δ_i+1} , $i = 1, 2$, then G_1 and G_2 pack.

This would complete the proof of Theorem 1.