

① Consider the following LP

$$\begin{aligned} \max \quad & 3x_1 + 4x_2 \\ \text{s.t.} \quad & 2x_1 + 5x_2 \leq 20 \\ & 4x_1 + 3x_2 \leq 24 \\ & x_1 + x_2 \geq 2 \\ & x_1, x_2 \geq 0 \end{aligned}$$

- (a) Sketch the feasible set. Identify all the corners of the feasible region (give their x_1, x_2 coords). Name them by A, B, C, ...
- (b) Write the LP in standard form.
- (c) What is the rank of the matrix A that you get?
- (d) Find the basic feasible solutions of this standard form LP that correspond to each corner point identified in part (a). You should clearly identify the basis B for each corner point and then solve $Bx_B = b$ to get the BFS corresponding to that B.
- (e) Do one iteration of the simplex method (5 steps described at the end of the class) starting from any one of the BFS found in (d).
- (f) Now add the fourth constraint $14x_1 + 7x_2 \leq 76$. Is there a degenerate BFS now? Why? If yes, then identify the corresponding corner point of the feasible region. Show that there are more than $n-m$ x_j 's set to zero at this basic solution.

②(a) Let x_j be a non-basic variable in a BFS x . Show that the reduced cost \bar{c}_j of x_j is

$$\bar{c}_j = c^T d$$

where d is j^{th} basic direction.

[Note x is a BFS of the LP]

$$\begin{array}{ll} \min & c^T x \\ \text{s.t.} & Ax = b \\ & x \geq 0 \end{array}$$

②(b) BT 3.2(a)

[Hint: you can give a very short proof using part (a). It would be useful to think of a direct proof as well.]

③ BT 3.5

Also, think about BT 3.1 and BT 3.2(b).

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This shows why we only need to find a locally optimal solution of a LP